

Optimal balanced circular packing problem

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Outline

- 1 Optimization problem and Applications
- 2 Nonlinear programming problem
- 3 Local search algorithm
- 4 Shor's dual bound ψ^*
- 5 Computational results

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Optimization problem

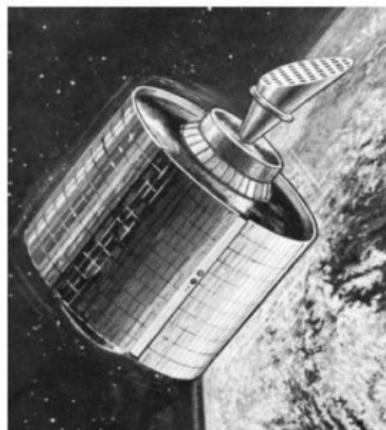
Let a collection of circles S_i , their radii r_i and weights w_i , $i = 1, \dots, m$, be given. We assume that the gravity centre of circle S_i is placed at its centre.

Optimization problem

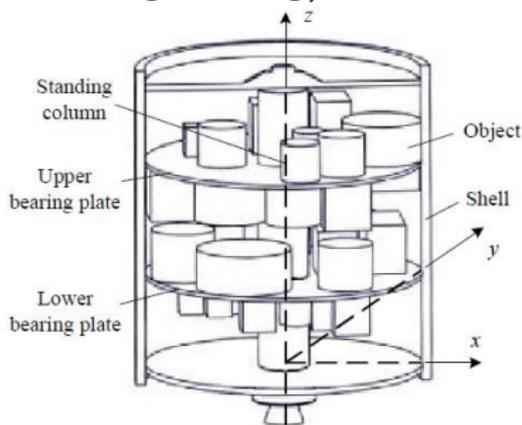
Pack the collection of circles S_i , $i = 1, \dots, m$, into containing circle S of minimal radius, provided that the gravity centre of the collection of circles is located at the centre of the circle S which coincides with the origin.

Application 1

for satellite modeling (space engineering)



a. the international communication satellite



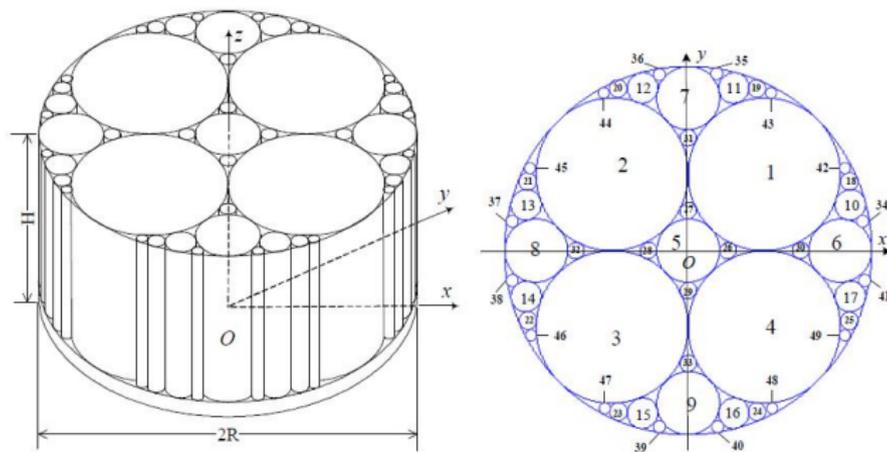
b. the simplified satellite module



FASANO G., PINTER J.D., EDS., Modeling and Optimization in Space Engineering. Springer Optimization and Its Applications. – Springer. – New York. – 2013. – 404 p.

Application 2

Cylinders and a container have the same height



CHE C., WANG Y., TENG H., Test problems for quasi-satellite packing: Cylinders packing with behavior constraints and all the optimal solutions known, Optimization Online (2008).

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Notation

Variables

We denote a variable centre of circle S_i by (x_i, y_i) and a variable radius of containing circle S by r .

Parameters

Let $\lambda_i = w_i / \sum_{i=1}^m w_i$, $i = 1, \dots, m$, and $r_{low} = \max_i r_i$ be the trivial lower bound of r .

Nonlinear programming problem

find

$$r^* = \min_{x,y,r} r \quad (1)$$

subject to

$$x_i^2 + y_i^2 \leq (r - r_i)^2, \quad i = 1, \dots, m, \quad (2)$$

$$(x_i - x_j)^2 + (y_i - y_j)^2 \geq (r_i + r_j)^2, \quad 1 \leq i < j \leq m, \quad (3)$$

$$\sum_{i=1}^m \lambda_i x_i = 0, \quad \sum_{i=1}^m \lambda_i y_i = 0, \quad (4)$$

$$r \geq r_{low}, \quad (5)$$

where $x = (x_1, \dots, x_m)$, $y = (y_1, \dots, y_m)$.

Constraints (2) - (5) mean

$$\varphi_i(x, y, r) = x_i^2 + y_i^2 - (r - r_i)^2 \leq 0, \quad i = 1, \dots, m, \quad (2')$$

i.e. circles S_i , $i = 1, \dots, m$ are contained into the containing circle S , if constraint (5) is satisfied

$$\phi_{ij}(x, y) = -(x_i - x_j)^2 - (y_i - y_j)^2 + (r_i + r_j)^2 \leq 0, \quad 1 \leq i < j \leq m, \quad (3')$$

i.e. each pair of circles from the collection S_i , $i = 1, \dots, m$ do not overlap,

$$\Phi_1(x, y) = \sum_{i=1}^m \lambda_i x_i = 0, \quad \Phi_2(x, y) = \sum_{i=1}^m \lambda_i y_i = 0, \quad (4')$$

i.e. gravity center of the collection of circles S_i , $i = 1, \dots, m$ is placed at the origin (the center of the containing circle S).

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Step 1: we reduce the problem (1)–(5)

to an unconstrained nonsmooth minimization problem

$$\min_{x \in R^m, y \in R^m, r} F(x, y, r), \quad (*)$$

where

$$F(x, y, r) = r + P_1 \left(\sum_{i=1}^m \max\{0, \varphi_i(x, y, r)\} + \sum_{i=1}^m \sum_{j>i}^m \max\{0, \phi_{ij}(x, y)\} \right) \\ + P_2 \sum_{k=1}^2 \max\{0, -\Phi_k(x, y) - \varepsilon_k, \Phi_k(x, y) - \varepsilon_k\} + P_3 \max\{0, -r + r_{low}\}.$$

Here P_1, P_2, P_3 are nonsmooth penalties (positive).

Step 2

We use the combination
of Shor's r -algorithm and multistart method
to find "good" local extrema
for the problem (*).

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The quadratic extremal problem

find

$$f^* = (r^*)^2 = \min_{r,x,y} r^2, \quad (6)$$

subject to

$$x_i^2 + y_i^2 - r^2 + 2r_i r - r_i^2 \leq 0, \quad i = 1, \dots, m, \quad (7)$$

$$-x_i^2 + 2x_i x_j - x_j^2 - y_i^2 + 2y_i y_j - y_j^2 + (r_i + r_j)^2 \leq 0, \quad 1 \leq i < j \leq m, \quad (8)$$

$$\sum_{i=1}^m \sum_{j=1}^m \lambda_i \lambda_j x_i x_j = 0, \quad \sum_{i=1}^m \sum_{j=1}^m \lambda_i \lambda_j y_i y_j = 0, \quad (9)$$

$$r^2 - (r_{low} + r_{up})r + r_{low}r_{up} \leq 0. \quad (10)$$

Here constraints (7) and (8) correspond to constraints (2) and (3), two equations in (9) correspond to equations (4). Quadratic inequality (10) follows from $r_{low} \leq r \leq r_{up}$, where r_{up} may be found by the local search algorithm.

Shor's dual bound ψ^*

gives a lower approximation of the objective function in the quadratic extremal problem.

The value of ψ^* can be found with given accuracy by methods of minimization of nonsmooth convex functions.



SHOR N.Z. Nondifferentiable optimization and polynomial problems. – Kluwer Academic Publishers (1998)

Properties of bound ψ^*

1) $\psi^* \leq f^* = (r^*)^2$ (Shor, 1985)

2) $r^* \geq \sqrt{\psi^*} \geq r_{low}$ (from problem (6)-(10))

It means, that bound $\sqrt{\psi^*}$ can be used to prove the global minimum for the problem (1)-(5).

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Program ralgb5

For computational experiments the program ralgb5 was used.
This program realizes
Shor's $r(\alpha)$ -algorithm with adaptive step-size regulation.

Example 1 ($m = 5$)

Circles have different radii and weights:

- $r_1 = 0.1, r_2 = 0.2, r_3 = 0.3, r_4 = 0.5, r_5 = 0.8$
- $w_1 = 0.0785, w_2 = 0.314, w_3 = 0.7065, w_4 = 1.9625, w_5 = 5.024$

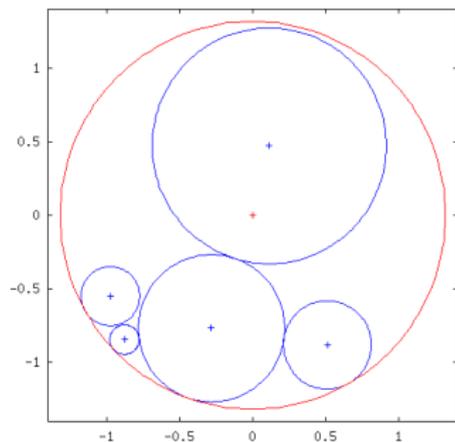
Tolerance for the components of the center of gravity

- $\varepsilon_1 = \varepsilon_2 = 0.0001$

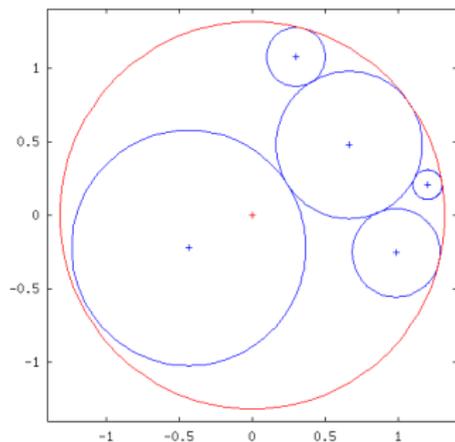


NENAKHOV E.I., ROMANOVA T.E., STETSYUK P.I.,
Balanced packing problem of circles in a circle of minimum
radius, Theory of optimal solutions, 143–153 (2013).

Two close local minima



$$r_{up} = 1.3175 \neq r^*$$



$$r_{up} = 1.3161 = r^* \text{ ???}$$

The proof that $r^* = 1.3161$

If

$$r_{low} = 0.8, \quad r_{up} = 1.35 > r^* = 1.3161$$

then

$$\psi^* = 1.7309$$

and we have

$$r^* \geq \sqrt{\psi^*} \geq 1.3156$$



STETSYUK P.I., ROMANOVA T.E., SCHIETHAUER G.
On the global minimum of the objective function in a balanced circular packing problem, *Dopovidi NAN Ukraine*, 6, 53–57 (2014).

Example 2 ($m = 7$)

Six circles have the same radius and weight:

- $r_1 = 1, r_2 = r_3 = r_4 = r_5 = r_6 = r_7 = 0.5$
- $m_1 = 4, m_2 = m_3 = m_4 = m_5 = m_6 = m_7 = 1$

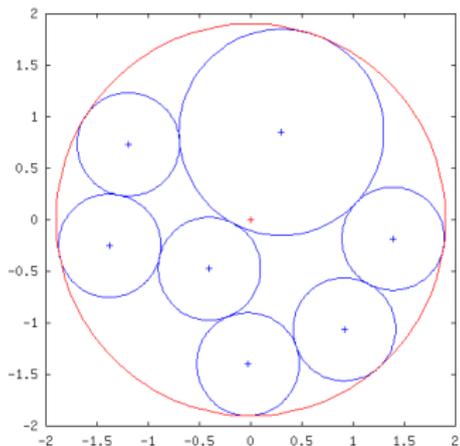
Tolerance for the components of the center of gravity

- $\varepsilon_1 = \varepsilon_2 = 0.0001$



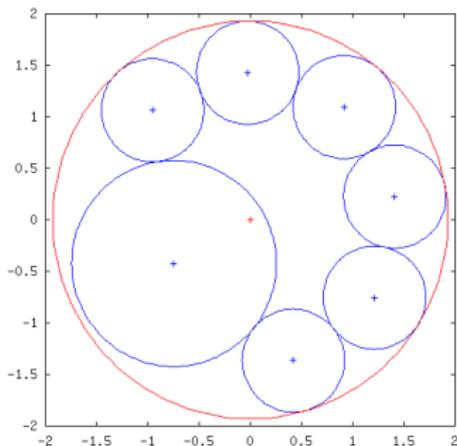
PSHENICHNYI B.N., SOBOLENKO L.A., Linearization method for inverse convex programming. Cybernetics and Systems Analysis, 31(6), 852-862 (1995).

Optimal placement of circles



without balance constraints

$$r_{up} = 1.9012 = r^*$$



with balance constraints

$$r_{up} = 1.9289 = r^* ???$$

There is no proof that $r^* = 1.9289$

If

$$r_{low} = 1.0, \quad r_{up} = 1.9290 > r^* = 1.9289$$

then

$$\psi^* = 2.4299$$

and we have

$$r^* \geq \sqrt{\psi^*} \geq 1.5588$$

Conclusion

Dual bound ψ^* can be successively improved by adding superfluous quadratic constraints to (7)–(10).

These constraints are non-trivial consequences of problem conditions. For instance, similar to (10) we can construct additional constraints for some of variables x_i, y_i , $i = 1, \dots, m$, using ranges of their variations.

Questions?

THANK YOU
FOR YOUR ATTENTION!

BACK UP SLIDES: Comments for program ralgb5

```

# Octave-code ralgb5 realizes Shor's  $r(\alpha)$ -algorithm
# with adaptive step for minimizing of convex function  $f(x)$ .
# ralgb5 uses user-defined octave's function  $[f,g]=\text{calcfg}(x)$ ,
# which calculate  $f=f(x)$  and subgradient  $g(x)$  at  $x$ .
# Input parameters:
#   calcfg - name of function for f and g calculation
#   x - starting point  $x_0(1:n)$  (is erased at the return)
#   alpha - the space dilation coefficient
#   h0, nh, q1, q2 - adaptive step parameters
#   epsx, epsg, maxitn - stop parameters
# Output parameters:
#   xr - the found minimum point  $x_r(1:n)$ 
#   fr - the function value at minimum point
#   itn - number of iterations
#   ncalls - number of calcfg calls
#   istop - the stop code (2=epsg,3=epsx,4=maxitn,5=error)

```

BACK UP SLIDES: Octave-function ralgb5

```

function [xr,fr,itn,ncalls,istop]=ralgb5(calcfg,x,alpha,h0,q1, # row001
                                     q2,nh,epsg,epsx,maxitn);
itn=0; hs=h0; B=eye(length(x)); xr=x; # row002
ncalls = 1; [fr,g0] = calcfg(xr); # row003
printf("itn %4d f %14.6e fr %14.6e ls %2d ncalls %4d\n", # row004
       itn, fr, fr, 0, ncalls);
if(norm(g0) < epsg) istop = 2; return; endif # row005
for (itn = 1:maxitn) # row006
    dx = B * (g1 = B' * g0)/norm(g1); # row007
    d = 1; ls = 0; ddx = 0; # row008
    while (d > 0) # row009
        x -= hs * dx; ddx += hs * norm(dx); # row010
        ncalls ++; [f, g1] = calcfg(x); # row011
        if (f < fr) fr = f; xr = x; endif # row012
        if(norm(g1) < epsg) istop = 2; return; endif # row013
        ls ++; (mod(ls,nh)==0) && (hs *= q2); # row014
        if(ls > 500) istop = 5; return; endif # row015
        d = dx' * g1; # row016
    endwhile # row017
    (ls == 1) && (hs *= q1); # row018
    printf("itn %4d f %14.6e fr %14.6e ls %2d ncalls %4d\n", # row019
          itn, f, fr, ls, ncalls);
    if(ddx < epsx) istop = 3; return; endif # row020
    xi = (dg = B' * (g1 - g0) )/norm(dg); # row021
    B += (1 / alpha - 1) * B * xi * xi'; # row022
    g0 = g1; # row023
endfor # row024
istop = 4; # row025
endfunction # row026

```

BACK UP SLIDES: The parameters for ralgb5

For minimization of nonsmooth functions we recommend:

$$\alpha = 2 \div 3, \quad h_0 = 1.0, \quad q_1 = 1.0, \quad q_2 = 1.1 \div 1.2, \quad n_h = 2 \div 3.$$

If the bound on distance from starting point x_0 to the minimum point x^* is given, then it is reasonable to choose initial step h_0 to be approximately equal to $\|x_0 - x^*\|$.

For minimization of smooth functions we recommend:

$$q_1 = 0.8 \div 0.95.$$

With this choice of parameters, as a rule, the number of descents in a direction rarely exceeds two, and for the n iterations the accuracy of the values of function is improved from three to five times.